

PySke: Algorithmic Skeletons for Python

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September 11, 2019

PySke

PySke = Python + Skeletons
Aims to easing parallelism

PySke: Python

Why Python? Because Python is cool! (But pythons are not).

- ▶ langage orienté objet simple à utiliser et à mettre en oeuvre
- ▶ Lambda calcul easy to use
- ▶ Typage dynamique, pas de declaration de variable (Smiley interrogatif) (typage possible depuis 3.7 mais pour cb de temps?)

Very used language in the programmer community

PySke: Python

Community graph 1 (Stack overflow researches)

PySke: Python

Community graph 2 (Google trend)

PySke: Skeleton

Definition (Murray Cole) SPMD vs Global view Inspiration
Fonctionnelle

Pattern de calcul. PySke : Sur des structures de donnees Object
(notation pointee) + functional (return of a new structure) Style

MURRAY COLE. Algorithmic Skeletons: Structured Management
of Parallel Com- putation. In *MIT Press* 1989.

PySke: Skeleton

Title : PySke 1.0

Papiers: These + HPCS 2019

JOLAN PHILIPPE AND FRÉDÉRIC LOULERGUE. PySke: Algorithmic skeletons for Python. In *International Conference on High Performance Computing and Simulation (HPCS)*. IEEE, 2019.

JOLAN PHILIPPE. Systematic Development of Efficient Programs on Parallel Data Structures. Master Thesis.

PySke: Skeleton

Etat de l'art (tableau) TARGET OF THE LANGUAGE: Lists +
Trees

PySke: Functionnal inspiration

Definition of lists (functional inspiration and Bird-Meertens Formalism).

List A : | **nil**: List A
 | **cons**: A \rightarrow List A \rightarrow List A

RICHARD S. BIRD. An introduction to the theory of lists. In *Logic of Programming and Calculi of Discrete Design*. NATO ASI Series F

PySke: Inspiration

Definition of lists (functional inspiration and Bird-Meertens Formalism).

$$\begin{array}{l} \text{List A :} \\ \quad | \text{ nil: List A} \\ \quad | \text{ cons: A} \rightarrow \text{List A} \rightarrow \text{List A} \end{array}$$

Notations:

- ▶ $[] = \text{nil}$
- ▶ $h :: t = \text{cons } h \ t$
- ▶ $a :: b :: c :: [] = [a; b; c]$

RICHARD S. BIRD. An introduction to the theory of lists. In *Logic of Programming and Calculi of Discrete Design*. NATO ASI Series F

PySke: Inspiration

Map: applies a function to every element of the structure.

$$\text{map} : (A \rightarrow B) \rightarrow \text{List } A \rightarrow \text{List } B$$

$$\text{map } f [x_1; \dots; x_n] = [f x_1; \dots; f x_n]$$

PySke: Inspiration

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Formally:

$$\begin{cases} \text{map } f [] & = [] \\ \text{map } f (x :: xs) & = (f x) :: (\text{map } f xs) \end{cases}$$

PySke: Inspiration

Reduce: collapse all the elements of the structure.

$$\text{reduce} : (A \rightarrow A \rightarrow A) \rightarrow \text{List } A \rightarrow A$$

$$\text{reduce } (\oplus) [x_1; \dots; x_n] = x_1 \oplus \dots \oplus x_n$$

PySke: Inspiration

Reduce: collapse all the elements of the structure.

$$\text{reduce} : (A \rightarrow A \rightarrow A) \rightarrow \text{List } A \rightarrow A$$

$$\text{reduce } (\oplus) [x_1; \dots; x_n] = x_1 \oplus \dots \oplus x_n$$

Formally:

$$\begin{cases} \text{reduce } \oplus [] & = \iota_{\oplus} \\ \text{reduce } \oplus (x :: xs) & = x \oplus (\text{reduce } \oplus xs) \end{cases}$$

PySke: Inspiration

Scan: Accumulate values through the list

$$\text{scan} : (A \rightarrow A \rightarrow A) \rightarrow \text{List } A \rightarrow \text{List } A$$
$$\text{scan } (\oplus) [x_1; x_2; \dots; x_n] = [x_1; x_1 \oplus x_2; \dots; x_1 \oplus x_2 \oplus \dots \oplus x_n]$$

PySke: Inspiration

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Formally? Boring definition.

PySke: Inspiration

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Formally? Boring definition.

But not that much if \oplus is symmetric:

$\text{scan } (\oplus) l = \text{reverse} \circ (\text{scan}' (\oplus) l) \circ \text{reverse}$ with

$$\begin{cases} \text{scan}' \oplus [] & = [] \\ \text{scan}' \oplus (x :: xs) & = (\text{reduce } \oplus (x :: xs)) :: (\text{scan}' \oplus xs) \end{cases}$$